

Information Systems Security

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Chapter 9

Public-Key Cryptography and RSA

Outline

- Principles of Public-Key Cryptosystems
- RSA Algorithm

Private-Key Cryptography

- traditional **private/secret/single key** cryptography uses **one** key
- shared by both sender and receiver
- if this key is disclosed communications are compromised
- also is **symmetric**, parties are equal
- hence does not protect sender from receiver forging a message & claiming is sent by sender

Public-Key Cryptography

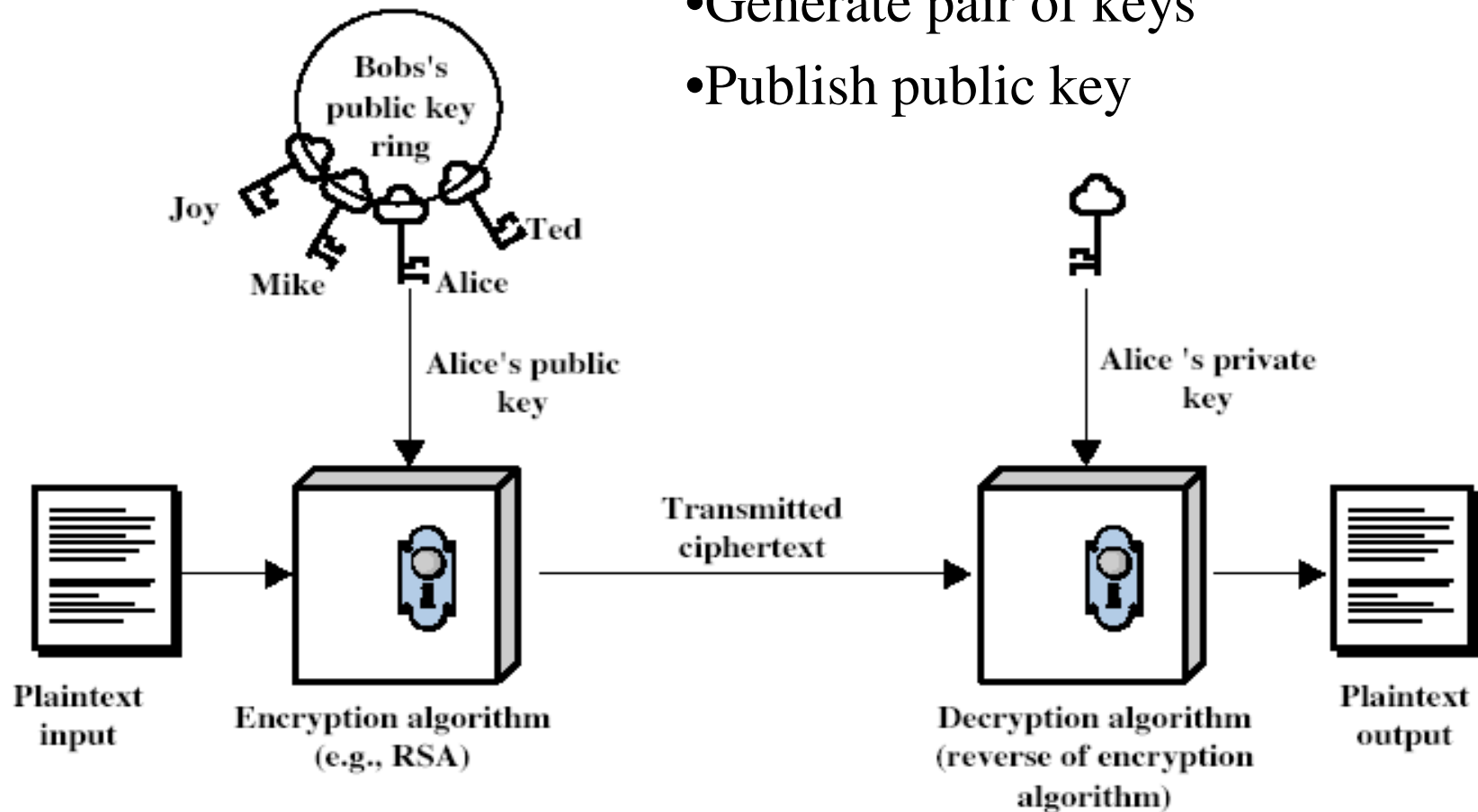
- probably most significant advance in the 3000 year history of cryptography
- uses **two** keys – a public & a private key
- **asymmetric** since parties are **not** equal
- uses clever application of number theoretic concepts to function
- complements **rather than** replaces private key crypto

Public-Key Cryptography

- **public-key/two-key/asymmetric** cryptography involves the use of **two** keys:
 - a **public-key**, which may be known by anybody, and can be used to **encrypt messages**, and **verify signatures**
 - a **private-key**, known only to the recipient, used to **decrypt messages**, and **sign** (create) **signatures**
- is **asymmetric** because
 - those who encrypt messages or verify signatures **cannot** decrypt messages or create signatures

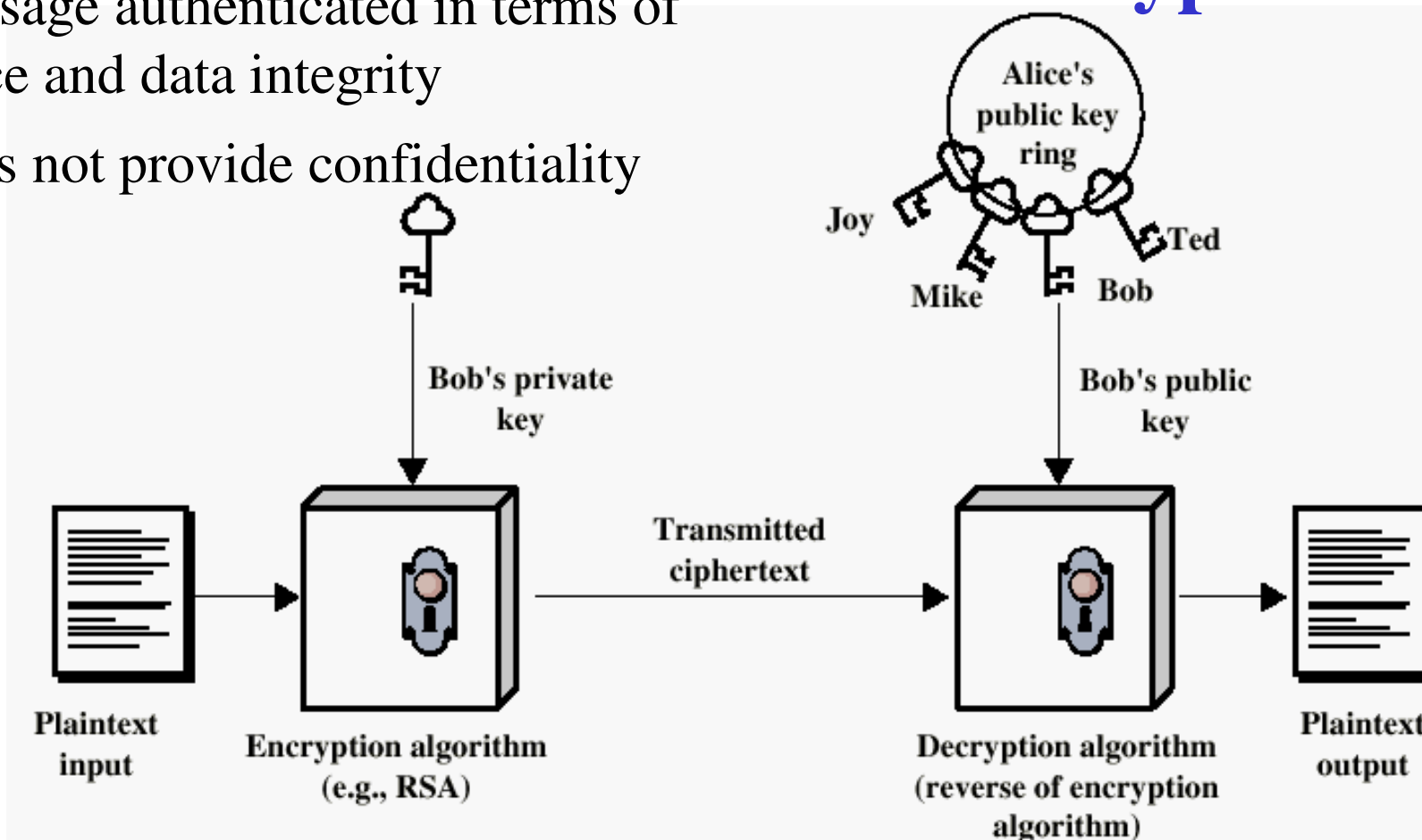
Public-Key Cryptography: Confidentiality

- Generate pair of keys
- Publish public key



- Entire encrypted message serves as a DS (can encrypt some bits as authenticator)
- Message authenticated in terms of source and data integrity
- Does not provide confidentiality

Authentication using Public-Key Crypto



Why Public-Key Cryptography?

- developed to address two key issues:
 - **key distribution** – how to have secure communications in general without having to trust a KDC with your key
 - **digital signatures** – how to verify a message comes intact from the claimed sender
- public invention due to *Whitfield Diffie & Martin Hellman* at Stanford Univ. in 1976
 - known earlier in classified community

Public-Key Characteristics

- Public-Key algorithms rely on two keys with the characteristics that it is:
 - computationally infeasible to find decryption key knowing only algorithm & encryption key
 - computationally easy to en/decrypt messages when the relevant (en/decrypt) key is known
 - either of the two related keys can be used for encryption, with the other used for decryption (in some schemes)

Public-Key Cryptosystems

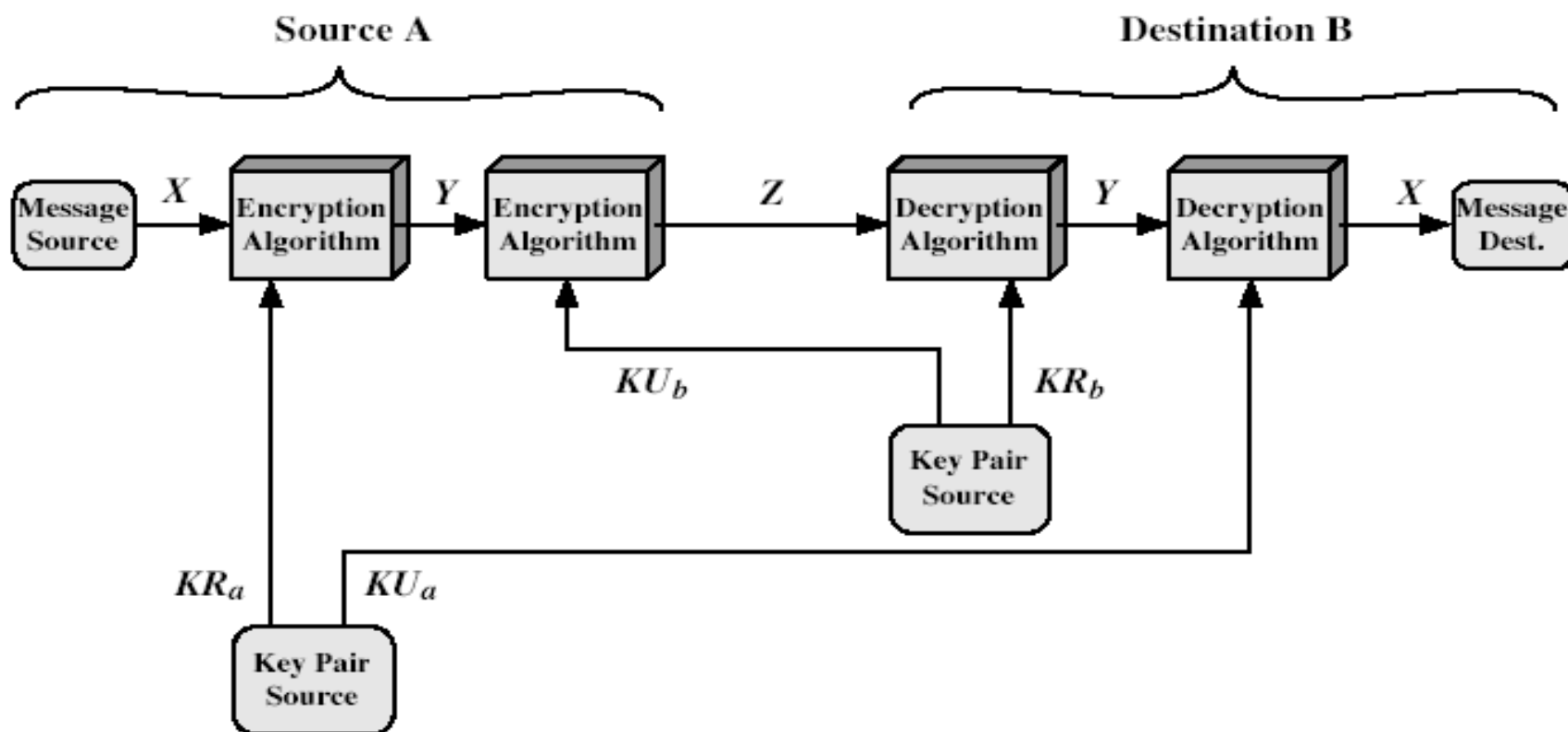


Figure 9.4 Public-Key Cryptosystem: Secrecy and Authentication

Public-Key Applications

- can classify uses into 3 categories:
 - **encryption/decryption** (provide secrecy → sender encrypts a message with the recipient's public key)
 - **digital signatures** (provide authentication → sender signs a message with its private key)
 - **key exchange** (of session keys)
- some algorithms are suitable for all uses, others are specific to one

Requirements for Public-Key Crypto

1. Computationally easy for a party B to generate a pair (public key KU_b , private key KR_b)

2. Easy for sender to generate ciphertext:

$$C = E_{KU_b}(M)$$

3. Easy for the receiver to decrypt ciphertext using private key:

$$M = D_{KR_b}(C) = D_{KR_b}[E_{KU_b}(M)]$$

Requirements for Public-Key Crypto

4. Computationally infeasible to determine private key (KR_b) knowing public key (KU_b)
5. Computationally infeasible to recover message M , knowing KU_b and ciphertext C
6. Encryption and decryptions functions can be applied in either order

$$M = D_{KRb}[E_{KU_b}(M)] = D_{KU_b}[E_{KRb}(M)]$$

Security of Public Key Schemes

- like private key schemes brute force **exhaustive search** attack is always theoretically possible
- **but keys used are too large (>512bits)**
- security relies on a **large enough** difference in difficulty between **easy** (en/decrypt) and **hard** (cryptanalysis) problems
- requires the use of **very large numbers**
- hence is **slow** compared to private key schemes
- *Public-key encryption currently confined to key management and signature applications*

RSA

- by Rivest, Shamir & Adleman of MIT in 1977
- best known & widely used public-key scheme
- Block cipher (use large numbers $n = 1024$ bits)
- For plaintext block M and ciphertext block C
 - $C = M^e \bmod n$
 - $M = C^d \bmod n$
 - Sender and receiver know n
 - Sender knows e
 - Receiver knows d
 - Public key $KU = \{e, n\}$
 - Private key $KR = \{d, n\}$

RSA Key Setup

- each user generates a public/private key pair by:
- selecting two large primes at random - p, q
- computing their system modulus $n=p \cdot q$ (factorization of large numbers)
 - note $\phi(n) = (p-1)(q-1)$
- selecting at random the encryption key e
 - where $1 < e < \phi(n), \text{gcd}(e, \phi(n)) = 1$
- solve following equation to find decryption key d
 - $e \cdot d = 1 \pmod{\phi(n)}$ and $0 \leq d \leq n$
- publish their public encryption key: $KU = \{e, n\}$
- keep secret private decryption key: $KR = \{d, p, q\}$

RSA Use

- to encrypt a message M , the sender:
 - obtains **public key** of recipient $KU = \{ e, n \}$
 - computes: $C = M^e \pmod n$, where $0 \leq M < n$
- to decrypt the ciphertext C , the receiver:
 - uses their private key $KR = \{ d, p, q \}$
 - computes: $M = C^d \pmod n$
- note that the message M must be smaller than the modulus n (block if needed)

RSA Example

1. Select primes: $p=17$ & $q=11$
2. Compute $n = pq = 17 \times 11 = 187$
3. Compute $\phi(n) = (p-1)(q-1) = 16 \times 10 = 160$
4. Select e : $\gcd(e, 160) = 1$; choose $e=7$
5. Determine d : $d \cdot e = 1 \pmod{160}$ and $d < 160$
Value is $d=23$ since $23 \times 7 = 161 = 1 \times 160 + 1$
6. Publish public key $KU = \{7, 187\}$
7. Keep secret private key $KR = \{23, 17, 11\}$

RSA Example cont.

- sample RSA encryption/decryption is:
- given message $M = 88$ (note that $88 < 187$)

- encryption:

$$C = 88^7 \bmod 187 = 11$$

- decryption:

$$M = 11^{23} \bmod 187 = 88$$